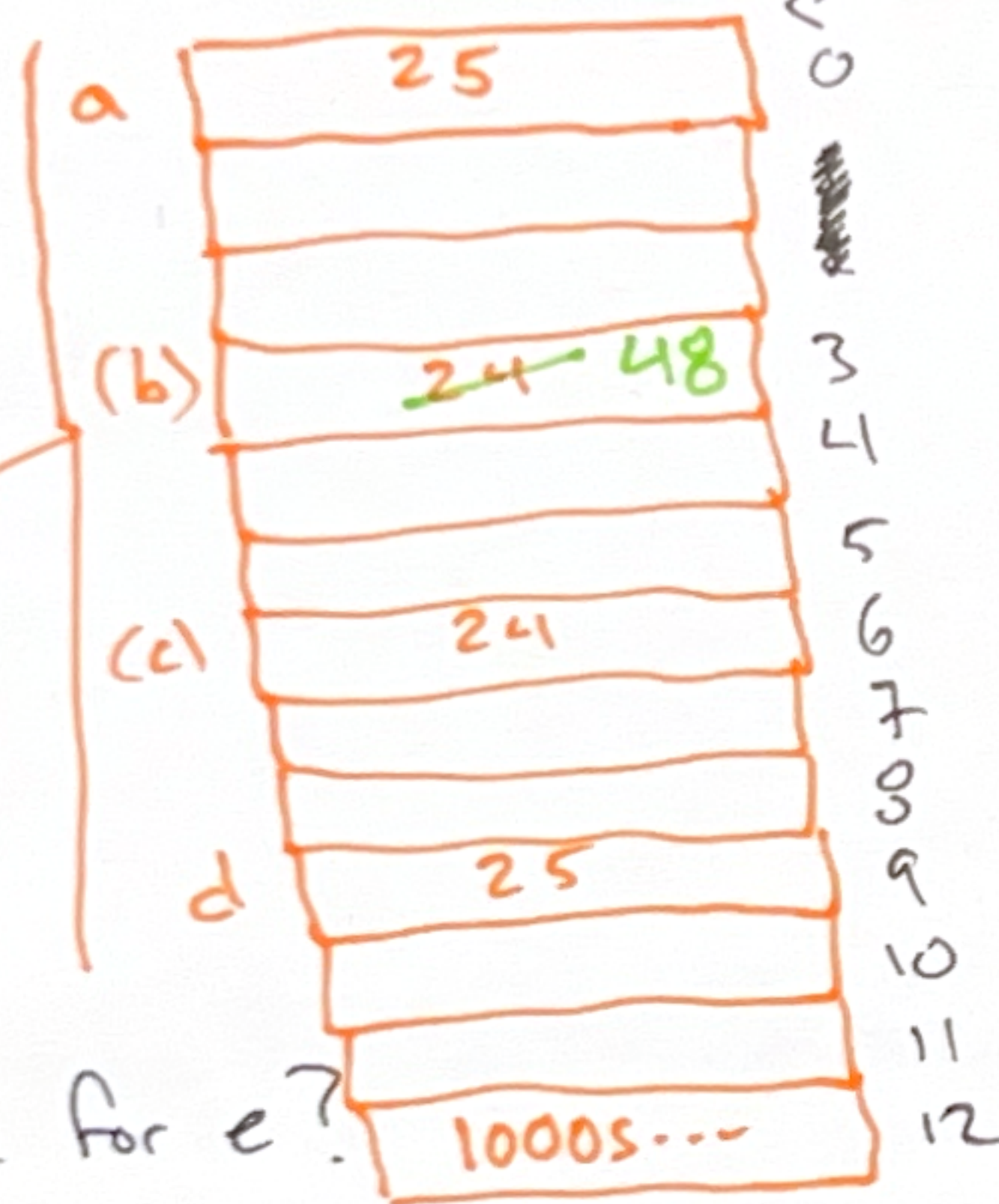


What about adjacent free blocks?

`a = malloc(16)`
`b = malloc(16)`
`c = malloc(16)`
`d = malloc(16)`
`free(b)`
`free(c)`
`e = malloc(32)`



i	start	header	next i	next h
0		25	3	24
3		24	6	24
6		24	9	25
9		25	12	1000....

Can we re-use b/c space for e?

Wednesday malloc would allocate at end of heap!

We need to coalesce adjacent free blocks.
 merge
 consolidate

Ideas?

- every time you malloc
in class

- on free, check next block too PA

```
size_t block_size_of(size_t user_request) {
```

```
}
```

```
uint64_t* find_block(uint64_t* start, size_t size) {
```

```
    int i = 0;
```

```
    while(i < HEAP_WORDS) {
```

```
        uint64_t header = start[i];
```

```
        if (header % 2 == 0 && nexth % 2 == 0) {
```

merging →

```
            start[i] = header + nexth;
```

```
            continue;
```

```
        } if (header % 2 == 0 && header == size)
```

```
            return &start[i];
```

```
        i += header / 8;
```

```
    }
```

```
    return NULL;
```

```
}
```

```
void split_block(uint64_t* block_start, size_t size) {
```

```
}
```

$(\text{header} \% 2) == 0$ // is the block free?

$\text{header} \geq \text{size}$ // what is its size?

insert
these
 $\text{int nexti} = i + \text{header} / 8$ // is the next block free?

$\text{uint64_t nexth} = \text{start}[\text{nexti}];$
 $(\text{nexth} \% 2) == 0$

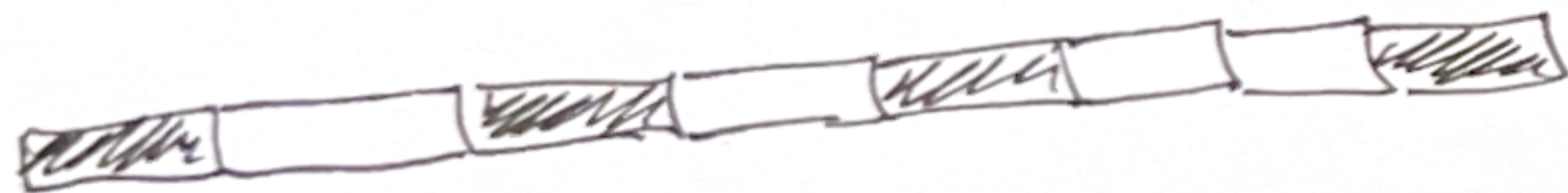
"Best fit" vs. "First fit" → use the earliest big enough block size

↳ look for closest big enough block size

"smallest big enough block"

Utilization vs. Speed → how many blocks are inspected per allocation

↳ what % of the heap can be used
or how much extra free space is between
start and last allocated block



↗
better utilization